

# Frequency-Domain Identification of Discrete-Time Systems using Sum-of-Rational Optimization

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Frequency-domain sysid is nonconvex

Can be cast as sum-of-rationals problem

Get global bounds on suboptimality using moment-SOS

# The Sysid Problem

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# System Identification Problem

Want to recover true SISO model  $(a, b) \in \mathbb{R}^{2n}$

$$G(z; a, b) = \frac{B(z; b)}{1 + A(z; a)} = \frac{\sum_{k=1}^n b_k z^{-k}}{1 + \sum_{k=1}^n a_k z^{-k}}$$

But we only have noisy data  $\mathcal{D} = \{(\omega_f, G_f), f = 1..N_f\}$

$$G_f = G_o(e^{j\omega_f}) + \eta_f$$

Optimize to minimize residuals  $\{\eta_f\}$

Least squares estimator for  $\theta = (a, b)$  given weights  $\{W_f\}$ :

$$\theta^* := \arg \min_{\theta \in \mathbb{K}} \mathcal{J}_{\mathcal{D}}(\theta)$$

$$\mathcal{J}_{\mathcal{D}}(\theta) = \sum_{f=1}^{N_f} \left| W_f (G_f - G(e^{j\omega_f}; \theta)) \right|^2$$

$W_f$  user-defined weights

$\mathbb{K} \subset \mathbb{R}^{2n}$  compact feasibility set

**Challenge:** global optimization over stable rational models!

- Classical methods: mainly *local search techniques*
  - + several heuristic initialization methods[Söderström & P. Stoica 1989], [Ljung 1999], [Pentilon & Schoukens 2012], ...
- *Numerical linear algebra* method  
e.g. [Agudelo et al. 2021], [Lagauw et al. 2023], ...
- Method based on the *Moments-SOS hierarchy*;  
e.g. [Rodrigues et al. 2019], [Rodrigues et al. 2020], [Vuillemin 2014]

Main contributions of this work:

- Use sum-of-rational structure
- infinite-dimensional LP, solved by hierarchy of *finite* SDPs.
- convergence in objective with increasing degree
- stability enforced using Polynomial Matrix Inequality (PMI)
- *Global optimally certified a-posteriori*: (rank)

# Sum-of-Rationals Reformulation

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## Sum-of-Rational Formulation: cost function

The SysID problem:  $P^* = \min_{\theta \in \mathbb{K}} \mathcal{J}_{\mathcal{D}}(\theta)$

$$\mathcal{J}_{\mathcal{D}}(\theta) = \sum_{f=1}^{N_f} \left| W_f (G_f - G(e^{j\omega_f}; \theta)) \right|^2, \quad G(z; \theta) = \frac{B(z; b)}{1 + A(z; a)}$$

## Sum-of-Rational Formulation: cost function

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Expand out each frequency data-point:

$$\begin{aligned} W_f (G_f - G(e^{j\omega_f}; \theta)) &= W_f \left( G_f - \frac{B(e^{-j\omega_f}; b)}{1 + A(e^{-j\omega_f}; a)} \right) \\ &= \frac{W_f \left( G_f (1 + A(e^{-j\omega_f}; a)) - B(e^{-j\omega_f}; b) \right)}{1 + A(e^{-j\omega_f}; a)} \end{aligned}$$

## Sum-of-Rational Formulation: cost function

Define polynomials  $p_f, q_f$  for each datapoint  $f$

$$\frac{p_f(\theta)}{q_f(a)} := \frac{\left| W_f \left( G_f(1 + A(e^{-j\omega_f}; a)) - B(e^{-j\omega_f}; b) \right) \right|^2}{\left| 1 + A(e^{-j\omega_f}; a) \right|^2}$$

$p_f(\theta)$  and  $q_f(a)$  **quadratic polynomials** in  $\theta$  and  $a$  respectively.

$$\sum_{f=1}^{N_f} \left| W_f(G_f - G(e^{j\omega_f}; \theta)) \right|^2 = \min_{\theta \in \mathbb{K}} \sum_{f=1}^N \frac{p_f(\theta)}{q_f(\theta)}$$

## Sum-of-Rational Formulation: feasibility set

Constraint term  $\mathbb{K} = \mathbb{K}_0 \cap \mathcal{S}_\delta$

- $\mathbb{K}_0$ : prior information on  $(a, b)$  range
- $\mathcal{S}_\delta$ : Enforce that  $(a, b)$  is a stable model

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$\mathcal{S}_\delta = \{\theta \in \mathbb{R}^{2n} : \Xi(a) \succeq \delta I\}$       Polynomial Matrix Inequality

with Hermite matrix  $\Xi(a) := \Theta(a)^T \Theta(a) - \tilde{\Theta}(a)^T \tilde{\Theta}(a)$ ,

$$\Theta(a) = \begin{bmatrix} 1 & a_1 & a_2 & \dots \\ 0 & 1 & a_1 & \dots \\ 0 & 0 & 1 & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix}, \quad \tilde{\Theta}(a) = \begin{bmatrix} a_n & a_{n-1} & a_{n-2} & \dots \\ 0 & a_n & a_{n-1} & \dots \\ 0 & 0 & a_n & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix} \in \mathbb{R}^{n \times n}$$

[Barnet 1983], [Henrion & Lasserre 2006]

## Wrapping up the formulation

$$\mathcal{J}_{\mathcal{D}} = \min_{\theta \in \mathbb{K}} \sum_{f=1}^N \frac{p_f(\theta)}{q_f(\theta)}$$

Objective from data  $\mathcal{D}$ , weights  $W$

Constraint set  $\mathbb{K}$  from prior info, stability

... now to solve it

# Bounds for Global Optimization

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# Transformation of Optimization

Original continuous optimization problem

$$f^* = \inf_{x \in C} f(x)$$

same objective as infinite-dimensional LP

$$f^* = \sup_{\gamma \in \mathbb{R}} \gamma$$
$$f(x) - \gamma \geq 0 \quad \forall x \in C$$

Hard problem: certification of nonnegativity over  $C$



# Ways to Discretize

Infinite-dimensional LP must be discretized for computation

More complexity: more accurate solutions

Method	Increasing Complexity
Gridding (MDP)	# Grid Points
Basis Functions (ADP)	# Functions
Random Sampling	# Samples
<b>Sum-of-Squares</b>	<b>Polynomial Degree</b>
Neural Nets (FOSSIL)	Width and Depth
Your Favorite Method	Some Accuracy Parameter

Runtime usually exponential in dimension, complexity

# Sum-of-Squares Tightening

Truncation when  $f, C$  have polynomial structure

$$f^* = \inf_{x \in C} f(x)$$

finite-degree  $d$  truncation

$$f_d^* = \sup_{\gamma \in \mathbb{R}} \gamma, \quad f(x) - \gamma \in \Sigma[C]_{\leq 2d}$$

Quadratic Module formed by constraint description

$$C = \{x \mid g_k(x) \geq 0 \forall k \in 0..N_c\}, \quad g_0(x) = 1$$

$$\Sigma[C] = \left\{ p \in \mathbb{R}[x] \mid p = \sum_{k=0}^{N_c} g_k(x) \sigma_k(x), \quad \deg(\sigma_k g_k) \leq 2d, \quad \sigma_k \in \Sigma[x] \right\}$$

# Sum-of-Rationals Tightening

SOS only admissible for polynomial structure.

Options for sum-of-rational optimization:

- Add new states
- Clear to common denominators
- Epigraph approach [Jibetean, de Klerk 2003]
- **Absolute-Continuity/Sandwich** [Bugarin, Henrion, Lasserre 2016]

# Sandwich approach

Sum-of-rationals optimization

$$P^* = \min_{\theta \in \mathbb{K}} \sum_{f=1}^{N_f} \frac{p_f(\theta)}{q_f(\theta)}$$

express as nonnegativity

$$= \max_{\gamma \in \mathbb{R}} \gamma, \quad \sum_{f=1}^{N_f} \frac{p_f(\theta)}{q_f(\theta)} - \gamma \geq 0 \quad \forall \theta \in \mathbb{K}$$

Assume all  $q_f(\theta) > 0$  over  $\mathbb{K}$  ( $e^{\pm j\omega_f}$  not a true pole)

# Sandwich approach

Introduce new functions  $\zeta \in \mathbb{R}[\theta]$ :

$$\frac{p_f(\theta)}{q_f(\theta)} \geq \zeta_f(\theta) \implies p_f(\theta) - \zeta_f(\theta)q_f(\theta) \geq 0$$

Substitute  $\zeta$  into objective

$$\max_{\gamma \in \mathbb{R}} \gamma, \quad \sum_{f=1}^{N_f} \zeta_f(\theta) - \gamma \geq 0 \quad \forall \theta \in \mathbb{K}$$

Now can truncate into degree- $\leq 2d$  SOS

# Recovery of Global Optima

SOS program is **dual** to LMI in moments

$$\left. \begin{aligned} P_d^* &= \min_{y, y^f} \sum_{f=1}^N \mathbb{L}_{y^f}(p_f(\theta)) \\ y_0 &= 1 \\ \mathbb{L}_{y^f}(\theta^\alpha q_f(a)) &= y_\alpha, \quad \forall \alpha \in \mathbb{N}_d^{2n} \\ \mathbb{M}_d(\mathbb{K}y) &\succeq 0, \quad \mathbb{M}_{d-1}(\mathbb{K}y^f) \succeq 0 \quad \forall f \end{aligned} \right\} \text{SDP}(d)$$

Read off candidate solutions from order-1 moments of  $\theta$

Global optima if  $\mathbb{M}_d(\mathbb{K}y)$ ,  $\mathbb{M}_{d-1}(\mathbb{K}y^f)$  satisfy rank conditions

# Computational Complexity and Convergence

Scaling linear in  $N_f$  and exponential in  $d$  (for fixed  $n$ )

Measure	Matrix	Size ( $\mathbb{K}_0$ )	Size ( $\mathbb{K}_s$ )	Mult.
$\mu$	$\mathbb{M}_d[\mathbb{K}y]$	$\binom{2n+d}{d}$	$n \binom{2n+d-1}{d-1}$	1
$\nu_f$	$\mathbb{M}_{d+1}[\mathbb{K}y^f]$	$\binom{2n+d+1}{d+1}$	$n \binom{2n+d}{d}$	$N_f$

Can exploit symmetry, term sparsity to reduce cost further

## Convergence Result:

- when  $\mathbb{K} = \mathbb{K}_0$ ,  $P_d^* \leq P_{d+1}^*$  for all  $d$
- when  $\mathbb{K} = \mathbb{K}_s$ ,  $P_d^* \uparrow P^*$

# Numerical Examples

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# Numerical Examples

MATLAB / julia

Data sets:

- $\omega_f = \frac{(f-1)\pi}{10}, \quad f = 1, \dots, 11$
- $G_f = G_o(e^{j\omega_f}) + \eta_f,$
- $\eta_f \sim \begin{cases} \mathcal{N}(0, 0.3^2), & f = 1, 11 \\ \mathcal{CN}(0, 0.3^2), & \text{otherwise} \end{cases}$



≡ ETFE with periodic input & transient effects ignored.

True systems with  $\forall f : W_f = 1$ :

- Case 1:  $G_o(z) = \frac{2z^{-1} - z^{-3}}{1 - 0.18z^{-1} - 0.134z^{-2} - 0.637z^{-3}}$
- Case 2: 14 randomly generated Schur-stable second-order systems

# Numerical Example 1

Stability is enforced:  $\mathbb{K}_S = \{\|\theta\|_\infty \leq 2\} \cap \mathcal{S}_{10^{-4}}$

Method	SDP(1)	n4sid	oe	fmincon(n4sid)	tfest
$\mathcal{J}_D(\theta^*)$	0.3173	0.3368	0.3224	0.3173	0.3173

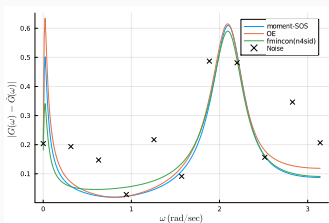
$\theta^*$  is extracted from the moment matrix  $\mathbb{M}_1[y^*]$

$$\theta_o = [-0.18 \quad -0.134 \quad -0.637 \quad 2 \quad 0 \quad -1 \quad ]^T$$

$$\theta^* = [-0.1415 \quad -0.2016 \quad -0.6107 \quad 1.978 \quad 0.0835 \quad -1.109]^T$$

$$0.3173 = \mathcal{J}_D(\theta^*) \geq P^* \geq P_1^* = 0.3173$$

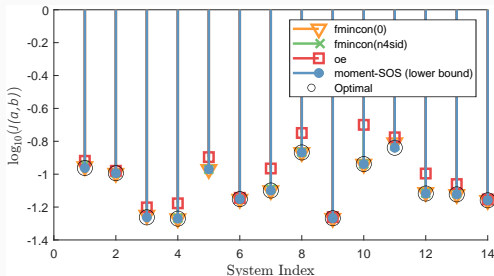
$$\implies P^* = P_1^*$$



## Numerical Example 2

Stability NOT enforced:  $\mathbb{K} = \{\|\theta\|_\infty \leq 2\} \leq 1.6$

- In all cases  $d = 3$  is used.
- almost all systems have global optimality certificate
- no solution is extracted for System 5



## Take-aways

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# Conclusions

An approach for frequency-domain identification of parameteric models

- tractable SDP + a certificate of global optimality
- stability enforced using Hermite criterion (PMI)

Future work includes:

- application to real systems (e.g. battery health)
- homogenization and improved conditioning
- comparisons with Sparse-POP , and similar methods.