

# Decomposed Structured Subsets for Semidefinite Optimization

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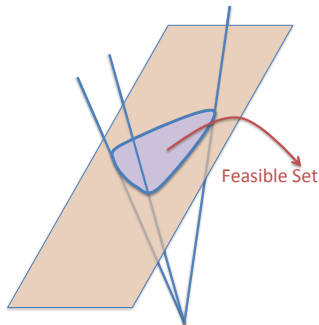
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# Semidefinite Programs

Modern control problems require solutions of large scale SDPs

$$\begin{aligned} \min_X \quad & \langle C, X \rangle \\ b_i = \quad & \langle A_i, X \rangle \quad |_{i=1}^m \\ X \in \quad & \mathbb{S}_+^n \end{aligned}$$

Scales polynomially with  $n$  and  $m$



# Approximations of SDPs

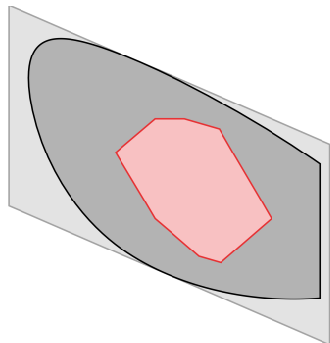
Approximate large SDPs  
Structured (simpler) subset of PSD

Diagonal Dominance ( $\mathcal{DD}$ ):

$$X_{ii} \geq \sum_{i \neq j} |X_{ij}|$$

bounds by Linear Programming

Generalizes to Factor Width



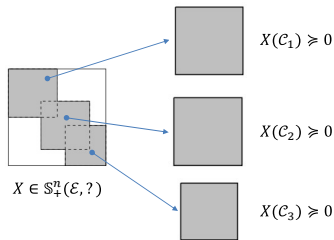
$$\mathcal{DD} \subset \mathcal{S}_+ \subset \mathcal{DD}^*$$

# Structure in SDPs

Use structure to simplify problem

- Sparsity
- Symmetry
- Algebraic

Large PSD cone  $\rightarrow$  smaller cones  
Often renders problems tractable



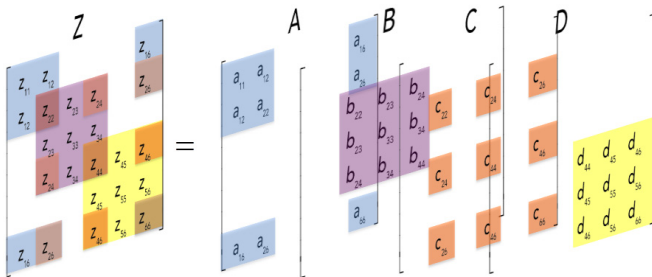
Blocks agree on overlap

# Sparse PSD Matrices

## Example of Chordal Sparsity

$$Z \succeq 0 \Leftrightarrow \exists A, B, C, D \text{ such that}$$

$$A + B + C + D = Z, \quad A, B, C, D \succeq 0$$



Large sparse PSD  $\leftrightarrow$  sum of smaller PSD

# Network H-infinity Estimation

Small Gain theorem:  $G$  robustly stable if  $\|G\|_\infty < 1$

Bounded Real Lemma:  $\|G\|_\infty < \gamma$  iff  $\exists P \succ 0$

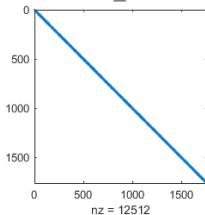
$$\begin{bmatrix} PA + A^T P + C^T C & P^T B + C^T D \\ B^T P + D^T C & -\gamma^2 I \end{bmatrix} \prec 0.$$

Block-diagonal  $P$  inherits sparsity of network

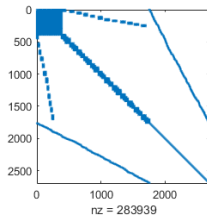
Sea Star Visualization



$P \geq \epsilon I$



Bounded Real Lemma



Sea Star network topology and LMI sparsity

# Sea Star

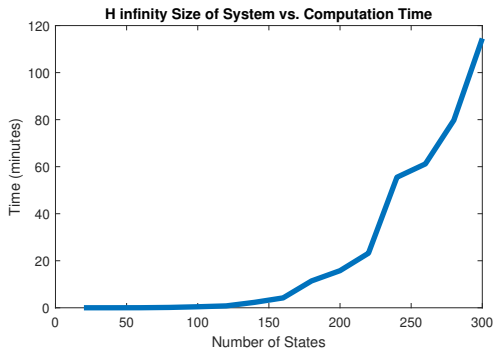


Inspiration for Sea Star Network

# H-infinity Optimization

No structure:  $\approx$  2 hours for 300 states

Out of memory:  $\geq$  400 states (64 GB RAM)



Dense  $H_\infty$  optimization

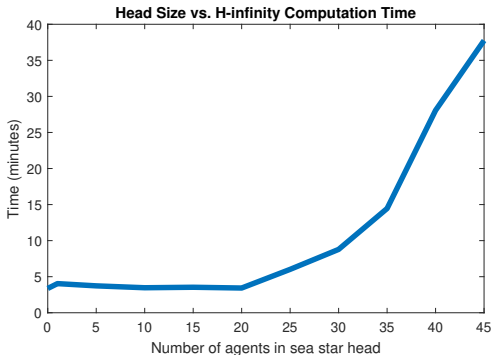


# H-infinity on Sea Star

Sea Star structure with chordal decomposition

Head sizes 0 . . . 45, largest clique 242

Out of memory: Head Size  $\geq 50$



Sparse  $H_\infty$  optimization

## Motivating Example $M(a,b)$

For what values of  $(a, b)$  is  $M(a, b)$  PSD-completable?

$$M(a, b) = \begin{bmatrix} 1 & \frac{1}{2} + a & ? & ? \\ \frac{1}{2} + a & 2 - b & -2a & a + b \\ ? & -2a & 5 & b/2 \\ ? & a + b & b/2 & 2 \end{bmatrix}$$

Cliques are  $\mathcal{C}_1 = (1, 2)$        $\mathcal{C}_2 = (2, 3, 4)$

# Decomposed Structured Subsets

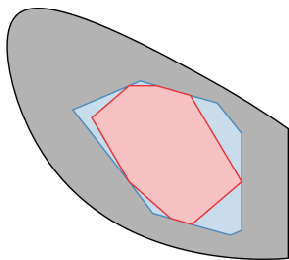
Structured Subsets destroy problem properties

Decompose, then approximate

Blocks are  $\mathcal{DD}$  vs. Matrix is  $\mathcal{DD}$

$$\begin{pmatrix} \boxed{X_{11} \quad X_{12}} & ? & ? \\ \boxed{X_{12} \quad X_{22}} & X_{23} & X_{24} \\ ? & X_{23} & X_{33} & X_{34} \\ ? & \boxed{X_{24} \quad X_{34}} & X_{34} & X_{44} \end{pmatrix} \in \mathbb{S}_+^4$$

Cliques:  $\{(1, 2), (2, 3, 4)\}$

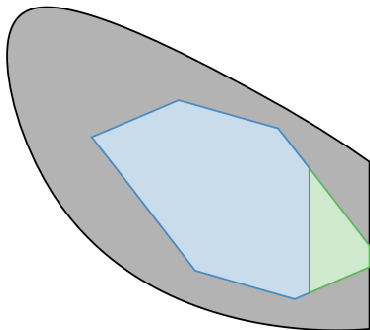


$$\mathcal{DD} \subset \{\mathcal{DD}, \mathcal{DD}\} \subset \mathbb{S}_+$$

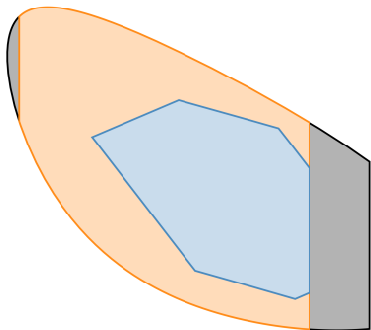
Preserving structure yields tighter approximations

# Mixing Cones

Adds flexibility in optimization



$\{S_+, DD\}$

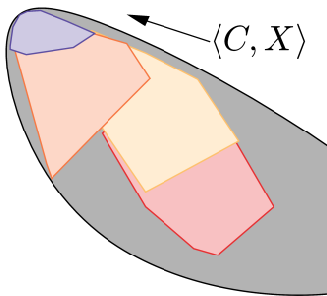


$\{DD, S_+\}$

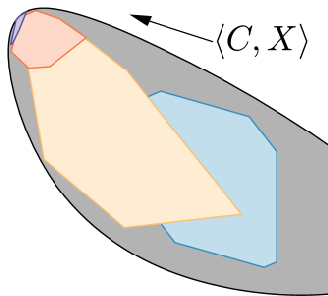
Useful if problem has few large cliques, bottleneck

# Change of Basis

Iteratively refine approximations (Ahmadi, Hall 2015)



Start from  $DD$



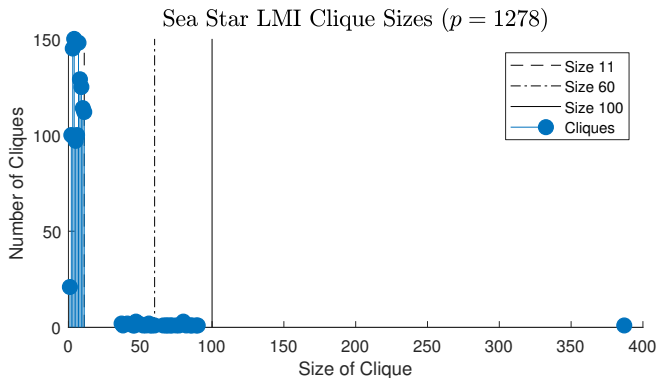
Start from  $\{DD, DD\}$

Standard vs. Decomposed Change of Basis

Improved solution quality with structure

# Sea Star Cliques

310 agents, each a random linear system of order  $\leq 10$   
70 in 'head', 20 in each of the 12 'arms' in 2 'knuckles'

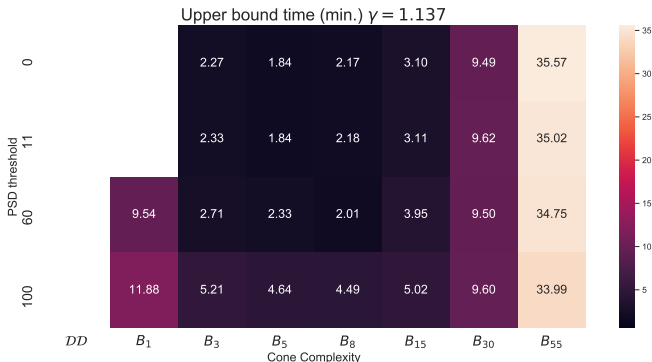


Single clique sized 387, rest have size  $< 90$

# Sea Star Upper Bounds

Cliques larger than threshold remain PSD

$B_k$ : Block-factor-width 2 matrices, k-sized blocks

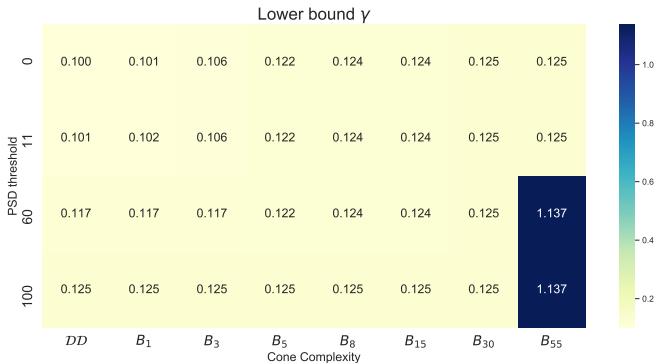


Time to find  $\gamma$  by upper bound  $K$  (minutes)

# Sea Star Lower Bounds

Dual decomposed structured subsets for lower bounds

Can verify if approximation = SDP solution (KKT)



$\gamma$  found by lower bound  $K^*$



# Conclusions

Structure in SDPs is powerful

Decompositions allow for tractable programs

Preserve structure while approximating

Increased performance and tighter bounds with structure

*<https://github.com/zhengy09/SDPfw>*